

## HW Assignment 2 (Due date: February 13, by 9:00 am)

1. [**Sorting Complexity, 5 points**] Exercise 1.2-2, page 14.
2. [**Time Complexity, 10 points**] Exercise 2.2-3, page 29.
3. [**Binary Search, 5 points**] Exercise 2.3-5, page 39.
4. [**Binary Search, 5 points**] Exercise 2.3-6, page 39.
5. [**Asymptotic Notation, 5 points**] Exercise 3.1-1, page 52.
6. [**Asymptotic Notation, 10 points**] Prove or disprove:
  - a)  $3^{n+1} = O(3^n)$
  - b)  $2^{2n} = O(2^n)$
  - c)  $3n^2 \lg n + 4n = O(n^3)$
  - d)  $3n^2 \lg n + 4n = O(n^2 \lg n)$
  - e)  $3n^2 \lg n + 4n = O(n^2 \sqrt{n})$
7. [**Asymptotic Notation, 5 points**] Prove  $\lg(n!) = \theta(n \lg n)$ .  
[Hint: use one of Stirling's approximations (3.18 or 3.20 on page 57)].
8. [**Substitution Method, 15 points**] Show that the solution to:
$$T(6) = 1$$
$$T(n) = 3T(\lfloor n/3 \rfloor + 4) + n, \text{ for } n > 6$$
is  $O(n \lg n)$ .
9. [**Master Method, 10 points**] Use the master method to give tight asymptotic bounds for the following recurrences:
  - a)  $T(n) = 4T(n/2) + n$ .
  - b)  $T(n) = 4T(n/2) + n^2$ .
  - c)  $T(n) = 4T(n/2) + n^3$ .
  - d)  $T(n) = 2T(n/4) + \sqrt{n}$ .
  - e)  $T(n) = 2T(n/4) + n^2$ .
10. [**Master Method, 5 points**] The recurrence  $T(n) = 10T(n/3) + n^2$  describes the running time of an algorithm  $A$ . A competing algorithm  $A'$  has a running time of  $T'(n) = aT'(n/9) + n^2$ . What is the largest integer value for  $a$  such that  $A'$  is asymptotically faster than  $A$ ?
11. [**Recurrence, 20 points**] Give asymptotic upper and lower bounds for  $T(n)$  in each of the following recurrences. Assume that  $T(n)$  is constant for  $n \leq 2$ . Make your bounds as tight as possible, and justify your answers.

- a)  $T(n) = 2T(n/2) + n^3$ .
- b)  $T(n) = T(9n/10) + n$ .
- c)  $T(n) = 16T(n/4) + n^2$ .
- d)  $T(n) = 7T(n/3) + n^2$ .
- e)  $T(n) = 7T(n/2) + n^2$ .
- f)  $T(n) = 2T(n/4) + \sqrt{n}$ .
- g)  $T(n) = T(n - 1) + n$ .
- h)  $T(n) = T(\sqrt{n}) + 1$ .

12. [**Recurrence, 5 points (\*)**]. Solve the recurrence  $T(n) = 2T(\sqrt{n}) + 1$  by making a change of variable. The solution should be asymptotically tight (i.e. use the  $\Theta$  notation). Do not worry about whether values are integral.
13. [**Design & Analysis, 10 points (\*)**] Exercise 2.3-7, page 39.