Optimal Stopping of Markov chain, Gittins Index and Related Optimization Problems

Isaac M. Sonin

Department of Mathematics and Statistics University of North Carolina at Charlotte, USA http://...Isaac Sonin in Google

New York, Columbia University, September 2011

- Optimal Stopping (OS) of Markov Chains
- State Elimination (SE) Algorithm
- *Gittins*, Katehakis Veinott and *w* (Whittle) Indices and their Generalizations
- Three Abstract Optimization Problems Abstract Optimization Equality
- Continue, Quit, Restart Probability Model
- Open Problems

There are two approaches - "Martingale theory of OS "and "Markovian approach".

Classical monographs:

- Chow, Robbins and Sigmund (1971)
- A. Shiryayev (1969, 1978)
- Dynkin, Yushkevich (1967)
- G. Peskir, A. Shiryayev (2006)
- T. Ferguson (website)

Optimal Stopping (OS) of Markov Chain (MC)

T. Ferguson: "Most problems of optimal stopping without some form of Markovian structure are essentially untractable...".

OS Model $M = (X, P, c, g, \beta)$: continue or stop

- X finite (countable) state space,
- $P = \{p(x, y)\}$, stochastic (transition) matrix
- c(x) one step cost function,
- g(x) terminal reward function,
- β discount factor, $0 \leq \beta \leq 1$
- (Z_n) MC from a family of MCs defined by a Markov Model M = (X, P)
- $v(x) = \sup_{\tau \ge 0} E_x[\sum_{i=0}^{\tau-1} \beta^i c(Z_i) + \beta^\tau g(Z_\tau)]$ value function

Description of OS Continues

- Remark ! absorbing state e, p(e, e) = 1, $p(x, y) \longrightarrow \beta p(x, y)$, $p(x, e) = 1 - \beta$, $\beta \longrightarrow \beta(x) = P_x(Z_1 \neq e)$ probability of "survival".
- $S = \{x : g(x) = v(x)\}$ optimal stopping set.

•
$$Pf = Pf(x) = \sum_{y} p(x, y)f(y)$$

Theorem (1, Shiryayev 1969)

(a) The value function v(x) is the minimal solution of Bellman equation ...

$$v = \max(g, c + Pv),$$

(b) if state space X is finite then set S is not empty and $\tau_0 = \min\{n \ge 0 : Z_n \in S\}$ is an optimal stopping time. ...

Basic methods of solving OS of MC, $c \equiv 0$

- The direct solution of the Bellman equation
- The value iteration method : one considers the sequence of functions $v_n(x) = \sup_{0 \le \tau \le n} E_x \dots, v_{n+1}(x) = max(g(x), Pv_n(x)), v_0(x) = g(x).$ Then $v_0(x) \le v_1(x) \le \dots v_n(x)$ converges to v(x).
- The linear programming approach ($|X| < \infty$), $\min \sum_{y \in X} v(y)$, $v(x) \ge \sum_{y} p(x, y)v(y), v(x) \ge g(x), x \in X$.
- Davis and Karatzas (1994), interesting interpretation of the Doob-Meyer decomposition of the Snell's envelope
- The State Elimination Algorithm (SEA) Sonin (1995, 1999, 2005, 2008, 2010)

State Elimination Algorithm for OS of MC

- OS = Bellman equation v(x) = max(g(x), c(x) + Pv(x));
- $M_1 = (X_1, P = P_1, c = c_1, g), S = S_1$. Three simple facts:
 - It may be *difficult* to find the states where it is optimal *to stop*, $g(x) \ge c_1(x) + P_1v(x)$, but it is *easy* to find a state (states) where it is optimal *not* to *stop*: *do not* stop if $g(z) < c_1(z) + P_1g(z)$ $\le c_1(z) + P_1v(z)$.
 - After identifying these states, set G, we can "eliminate" the subset
 D ⊂ G, and recalculate $P_1 → P_2$ and $c_1 → c_2$, g. Elimination
 theorem: $S_1 = S_2$, $v_1 = v_2$. Repeat these steps until
 $g(x) ≥ c_k(x) + P_k g(x)$ for all remaining $x ∈ X_k$. Then
 - Proposition 1. Let M = (X, P, c, g) be an optimal stopping problem, and g(x) ≥ c(x) + Pg(x) for all x ∈ X. Then X is the optimal stopping set in the problem M, and v(x) = g(x) for all x ∈ X.

Eliminate state(s) z, (set D) and recalculate probabilities

Embedded Markov chain (Kolmogorov, Doeblin) $M_1 = (X_1, P_1), D \subset X_1, X_2 = X_1 \setminus D, (Z_n) \text{ MC } \tau_0, \tau_1, ..., \tau_n, ..., \text{ the moments of zero, first, and so on, visits of <math>(Z_n)$ to the set X_2 . Let $Y_n = Z_{\tau_n}, n = 0, 1, 2, ...$

Lemma (KD)

(a) The random sequence (Y_n) is a Markov chain in a model $M_2 = (X_2, P_2)$, where $P_2 = \{p_2(x, y)\}$ is given by formula $P_1 = \begin{bmatrix} Q & T \\ R & P_0 \end{bmatrix}$, $P_2 = P_0 + RU = P_0 + RNT$, N is a (transient) fundamental matrix, i.e. $N = (I - Q)^{-1}$, $N = I + Q + Q^2 + ... = (I - Q)^{-1}$, U = NT.

State Reduction Approach: GTH/S algorithm to calculate the invariant distribution (1985) for $D = \{z\}$

State Elimination Algorithm, $c \equiv 0$

If $D = \{z\}$ then

$$p_2(x,y) = p_1(x,y) + p_1(x,z)n_1(z)p_1(z,y),$$

where $n_1(z) = 1/(1 - p_1(z, z))$.

State Reduction approach

GTH/S algorithm (1985), invariant distr. for ergodic MC; W. Grassmann, M. Taksar, D. Heyman, (1985), T. J. Sheskin (1985).

State Elimination for OS of MCs

Presman, E., Sonin, I. (1972). The problem of best choice ... a random number of objects.

I. M. Sonin, (1995,1999). The Elimination Algorithm ...Math. Meth. of Oper. Res., Advances in Mathematics

Irle, A. (1980). On the Best Choice Problem with Random Population Size. Z.O.R., MC (2006)

E. Presman (2011) Continuous time OS

There is a well known connection between three problems related to Optimal Stopping of Markov Chain and the equality of three corresponding indices: the *classical Gittins index* in the *Ratio Maximization Problem*, the Kathehakis-Veinot index in a *Restart Problem*, and w (Whittle) index in a *family of Retirement Problems*.

from these indices \rightarrow to generalized indices \rightarrow to \ldots

One of the goals of my talk is to demonstrate that the equality of these (generalized) indices is a special case of a more general relation between three simple *abstract optimization* problems.

There is no doubt that the relationship between these problems was used in optimization theory before on different occasions in *specific problems* but we fail to find a *general statement* of this kind in the vast literature on optimization. **Reward Model** $M = (X, P, c(x), \beta)$, continue or stop.

Given a reward model M and point $x \in X$, the classical Gittins index, $\gamma(x)$, is defined as the maximum of the expected discounted total reward during the interval $[0, \tau)$ per unit of expected discounted time for the Markov chain starting from x, i.e.

$$\gamma(\mathbf{x}) = \sup_{\tau > 0} \frac{E_{\mathbf{x}} \sum_{n=0}^{\tau-1} \beta^n c(Z_n)}{E_{\mathbf{x}} \sum_{n=0}^{\tau-1} \beta^n}, \quad 0 < \beta = \text{const} \le 1.$$

Multi-armed bandit (MAB) Problems: a number of competing projects, each returning a stochastic reward. Projects are *independent* from each other and only one project at time may evolve. *Gittins Theorem*: Gittins index policy is optimal. Not true for *dependent* arms ! Classical case (D. Feldman, 1962). Presman, Sonin book (AP, 1990) on MAB problems. Bank, P; Follmer, H., American options, multi-armed bandits, and optimal consumption plans: a unifying view. Lecture Notes in Math., 1814, Springer, Berlin, 2003.

Bank, P; El Karoui, N., A stochastic representation theorem with applications to optimization and obstacle problems. Ann. Probab. 32 (2004), no. 1B, 1030-1067.

Bank, P; Kuchler, C., On Gittins' index theorem in continuous time. Stochastic Process. Appl. 117 (2007), no. 9, 1357-1371.

Gittins J., Glazebrook K., Weber R., Multi-armed Bandit Allocaton Indices, 2nd edition, Wiley, 2011.

KV index

KV index $M = (X, s, P, c(x), \beta)$, continue or restart to s.

Let h(x|s) denote the supremum over all strategies of the expected total discounted reward on the infinite time interval in reward model with an initial point x, and restart point s. Using the standard results of Markov Decision Processes theory, Katehakis and Veinott proved that function h(x|s) satisfies the equality

$$h(x|s) = \sup_{\tau>0} E_x[\sum_{n=0}^{\tau-1} \beta^n c(Z_n) + \beta^\tau h(s)],$$

and $\gamma(s) = (1 - \beta)h(s)$, where by definition h(s) = h(s|s). We call index h(s) a KV index. This index can be defined for any point $x \in X$, so we use also notation h(x).

w (Whittle) index

Retirement Process formulation was provided by Whittle (1980). Given a reward model M, he introduced the parametric family of OS models $M(k) = (X, P, c(x), k, \beta)$, where parameter k is a real number and the terminal reward function g(x) = k for all $x \in X$. Denote v(x, k) the value function for such a model, i.e. $v(x, k) = \sup_{\tau \ge 0} E_x[\sum_{n=0}^{\tau-1} \beta^n c(Z_n) + \beta^{\tau} k]$, and denote Whittle index

 $w(x) = \inf\{k : v(x,k) = k\}.$

Since $\beta < 1$, for sufficiently large k it is optimal to stop immediately and v(x, k) = k. Thus $w(x) < \infty$. The results of Whittle imply that v(x, k) = k for $k \ge w(x)$, v(x, k) > k for k < w(x), and w(x) = h(x).

Theorem 2

The three indices defined for a reward model $M = (X, P, c(x), \beta), 0 < \beta < 1$, coincide, i.e. $h(x) = w(x) = \gamma(x)/(1 - \beta), x \in X$.

Sonin (Stat. & Prob. Let., 2008): simple and transparent algorithm to calculate this common index. This algorithm is based on State Elimination algorithm. Nino-Mora (2007) for classical GI.

To apply this algorithm it is necessary to replace a constant discount factor β by a variable "survival" probability $\beta(x)$, because after the first recursive step a discount factor is not a constant anymore. So by necessity a more general model was considered and the classical GI $\gamma(x)$ was replaced by a generalized Gittins Index (GGI) $\alpha(x)$ as follows.

Generalized Gittins index

In general case, when $\beta(x)$ can be variable, we denote $P_x(Z_\tau = e)$ by $Q^{\tau}(x)$, the probability of termination on $[0, \tau)$, and we define the Generalized GI (GGI), $\alpha(x)$, for a model with termination as

$$\alpha(x) = \sup_{\tau > 0} \frac{R^{\tau}(x)}{Q^{\tau}(x)},$$

i.e. $\alpha(x)$ is the maximum discounted total reward per chance of termination.

Mitten, L. G. (1960); Denardo, Eric V.; Rothblum, Uriel G.; Van der Heyden, Ludo (2004) Index policies for stochastic search in a forest ... Math. Oper. Res.

The common part of all three problems described above is a maximization over set of all positive stopping times τ .

Maximization over the same set !

Three abstract indices

Three abstract optimization problems

Suppose there is an *abstract index set U*, and $A = \{a_u\}$ and $B = \{b_u\}$ be two sets indexed by the elements of *U*. Suppose that an assumption *U* holds,

 $a_u \leq a < \infty, \quad 0 < b \leq b_u \leq 1$ (U)

Problem 1. Restart problem (from Katehakis-Veinott index) Find

solution(s) of the equation

$$h = sup_{u \in U}[a_u + (1 - b_u)h])$$
, i.e.

$$h=H(h), \quad (*)$$

where $H(k) = \sup_{u \in U} [a_u + (1 - b_u)k]).$

h = Abstract KV index

There are two equivalent interpretations of this problem.

There is a set of "buttons" $u \in U$. A DM can select one of them and push. She obtains a reward a_u and according to the first interpretation with probability b_u the game is *terminated*, and with complimentary probability $1 - b_u$ she can select any button again. Her goal is to maximize the total (undiscounted) reward.

According to the second interpretation the game is continued sequentially and $1 - b_u$ is not the probability but a *discount factor* applied to the future rewards. It can be easily proved that in both cases its value satisfies the equation above.

Our second optimization problem is

Problem 2. Ratio (cycle) problem

Find α

 $\alpha = \sup_{u \in U} \frac{a_u}{b_u} \qquad \text{(Gittins index)} \quad (**)$

The interpretation of this problem is straightforward: a DM wants to maximize the ratio of the one step reward per "chance of termination".

 $\alpha = \text{Abstract Gittins index}$

Problem 3. A parametric family of Retirement problems

Find w, (abstract Whittle index) defined as follows: given parameter $k, -\infty < k < \infty$, let $v(k) = \max(k, H(k))$, where $H(k) = \sup_{u \in U} [a_u + (1 - b_u)k]).$ $w = \inf\{k : v(k) = k\}.$ (* * *) $h = \sup_{u \in U} [a_u + (1 - b_u)h]).$ (*) $\alpha = \sup_{u \in U} \frac{a_u}{b_u}$ (**)

Theorem 3

a) The solution h of the equation (*) exists, is unique and finite; b) $h = \alpha = w$; c) the optimal index u (or an optimizing sequence u_n) for any of three problems is the optimal index (or an optimizing sequence) for two other problems.

The proof is elementary. Function H(k) is nondecreasing, continuous, and convex.

Theorem 2 from Theorem 3: $U = \{u\} = \{\text{set of all Markov moments} \\ \tau > 0\}, a_u = R^{\tau}(x) = E_x \sum_{n=0}^{\tau-1} c(Z_n), \text{ the total expected reward till moment } \tau, \text{ and } b_u = Q^{\tau}(x) = P_x(Z_{\tau} = e), \text{ the probability of termination on } [0, \tau).$

Remark ! Theorem 3: Only equivalence, not how to solve !

Problem 4. Suppose that a DM has to solve the optimization problem similar to Problem 1 with sequential selection of buttons with only one distinction - every button can be used *at most once*.

The Mitten's result (1960) essentially can be described as

Theorem 4

Suppose that there is a sequence of indices u_n such that after the relabeling of indices in this sequence, we have $\alpha_1 = \frac{a_1}{b_2} \ge \alpha_2 = \frac{a_2}{b_2} \ge \ldots \ge \frac{a_u}{b_u}$ for each $u \in U$ not in this sequence. Then to push buttons in the order 1, 2, ... is an optimal strategy.

(joint with S. Steinberg) A general CQR model is specified by a tuple $M = (X, B, P, A(x), c, q, r_i(x))$, where X is a countable state space, $B = \{s_1, ..., s_m\}$ is a subset of a state space $X, P = \{p(x, y)\}$ a stochastic matrix

set of available actions $A(x) = (\text{continue}, \text{quit}, \text{restart to } s_i)$ reward function r(x|a) is specified by particular functions c(x), q(x) and $r_i(x), i = 1, 2, ..., m$.

If an action c, "continue" is selected then r(x|c) = c(x) and transition to a new state occurs according to transition probabilities p(x, y), if an action q, "quit" is selected then r(x|q) = q(x) and transition to an absorbing state e occurs with probability one, if an action r_i , "restart to state s_i " is selected then $r(x|r_i) = r_i(x)$ and transition to a state s_i occurs with probability one. *Idea:* Sonin (2008), St.& Pr. Let. no quit, restart with zero fee; a family of Whittle OS models $M(k) = (X, P, c(x), g(x|k), \beta(x))$, terminal reward function g(x|k) = k.

Consider function G(x|k) = g(x|k) - [c(x) + Pg(x|k)]. Linear function. Move k from large values to the left. Eliminate and Eliminate, order of n^3 .

Now, more complicated: g(x|k) = max(q(x), (r(x) + k)), G(x|k) is only piecewise linear ! Eliminate and Insert. Eliminate and Insert...

Theorem 5

For m = 1 and finite X there is an algorithm solving this problem in a finite number of steps; order of n^4 .

Open Problems.

- multiple restart problems
- multidimensional equivalent abstract optimization problems

 $\mathbf{h} = \sup_{u} [\mathbf{a}_{u} + \mathbf{B}_{u}\mathbf{h}], \text{ vectors and matrices or...}$

explanation of world financial crisis

Thank you for your attention !