

# Computing the Closeness Centrality of Evolving Networks on Clusters

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## Introduction

- Centrality metrics answer questions such as:
  - Who controls the flow in a network?
  - Who is more important? Who has more influence?
  - Whose contribution is significant for connections?

They have various applications:

- Covert network (e.g., terrorist identification).
- Contingency analysis (e.g., weakness/robustness of networks).
- Viral marketing. Traffic analysis. Store locations.

How to maintain an accurate analysis upon graph evolutions?

## Closeness Centrality

Let  $G = (V, E)$  be an unweighted graph with the vertex set  $V$  and edge set  $E$ .

The farness of a vertex is  $far[u] = \sum_v d(u, v)$ .

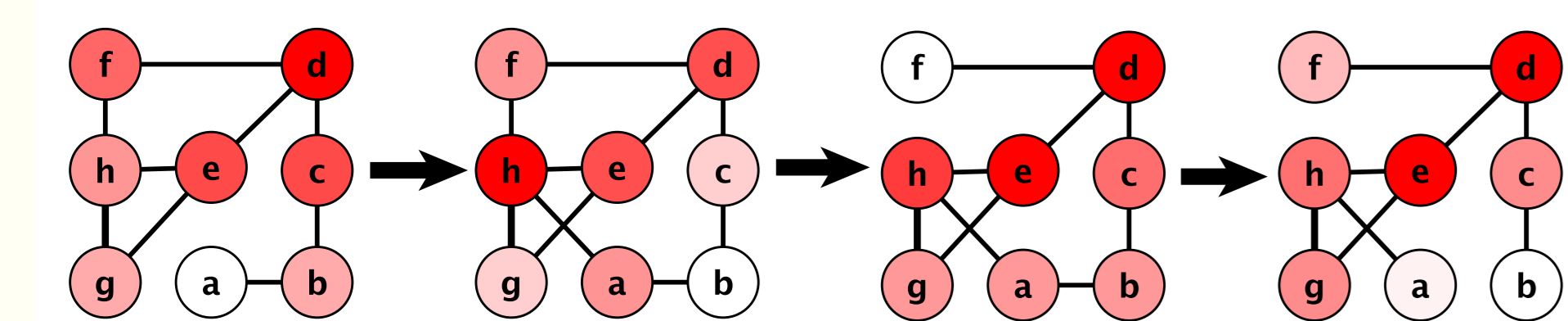
The closeness of a vertex is  $cc[u] = \frac{1}{far[u]}$ .

$d(u, v)$  is the length of the shortest path  $u \rightsquigarrow v$ .

The best known algorithm computes the shortest path graph rooted in each vertex of the graph.

The complexity is  $O(E)$  per source,  $O(VE)$  in total, which makes its computationally expensive.

## Variation of Closeness Centrality



The redness of a vertex represents its Closeness.

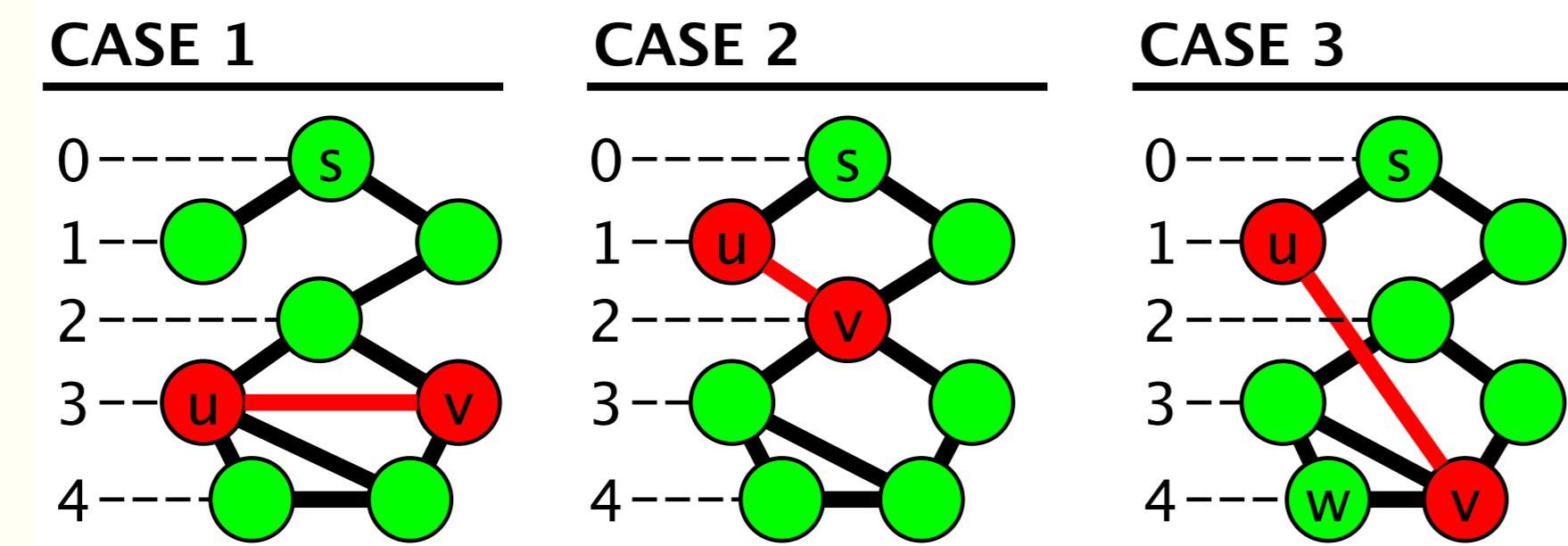
## Conclusion

Structural properties of Closeness Centrality lead to an efficient incremental algorithm. The algorithm scales on clusters and enables real-time Centrality analysis.

## References

- [1] A. E. Sariyüce, E. Saule, K. Kaya, and Ü. V. Çatalyürek. Hardware/software vectorization for closeness centrality on multi-/many-core architectures. In *Workshop on Multithreaded Architectures and Applications (MTAAP)*, 2014.
- [2] A. E. Sariyüce, E. Saule, K. Kaya, and Ü. V. Çatalyürek. Streamer: a distributed framework for incremental closeness centrality computation. In *IEEE Cluster*, 2013.
- [3] A. E. Sariyüce, E. Saule, K. Kaya, and Ü. V. Çatalyürek. Incremental algorithms for closeness centrality. In *IEEE BigData*, 2013.
- [4] A. E. Sariyüce, E. Saule, K. Kaya, and Ü. V. Çatalyürek. Shattering and compressing networks for betweenness centrality. In *SIAM International Conference on Data Mining, SDM*, 2013.

## Level Filtering

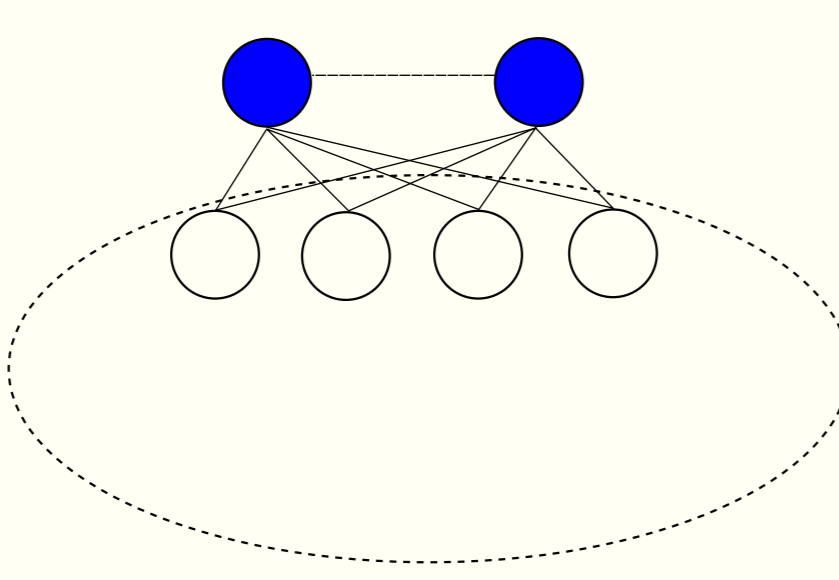


Case 1: No change.

Case 2: Only the structure changes.

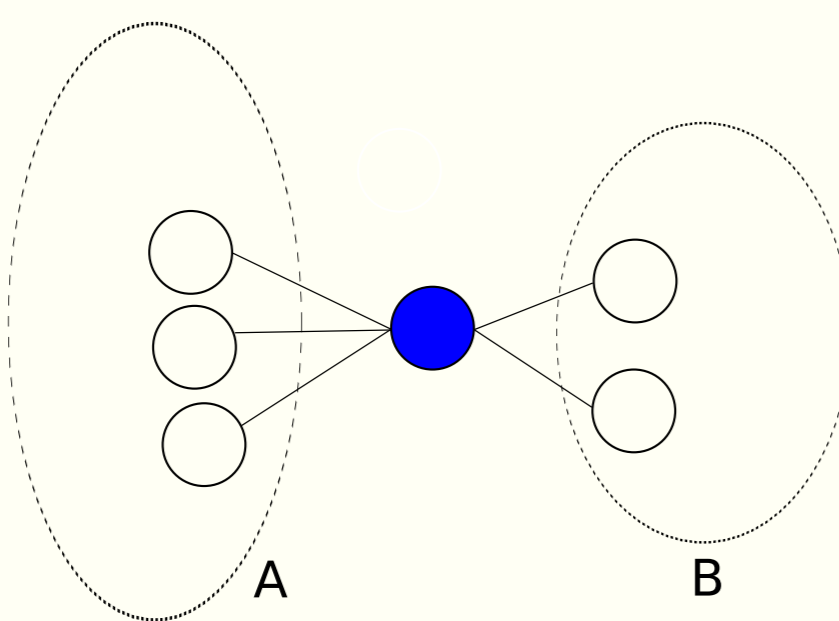
Case 3: Some distances change.

## Identical Vertices



Vertices identical to each other have the same farness.

## Articulation point



When an edge changes in  $A$ , all the vertices in  $B$  are impacted in the same way as the articulation point.

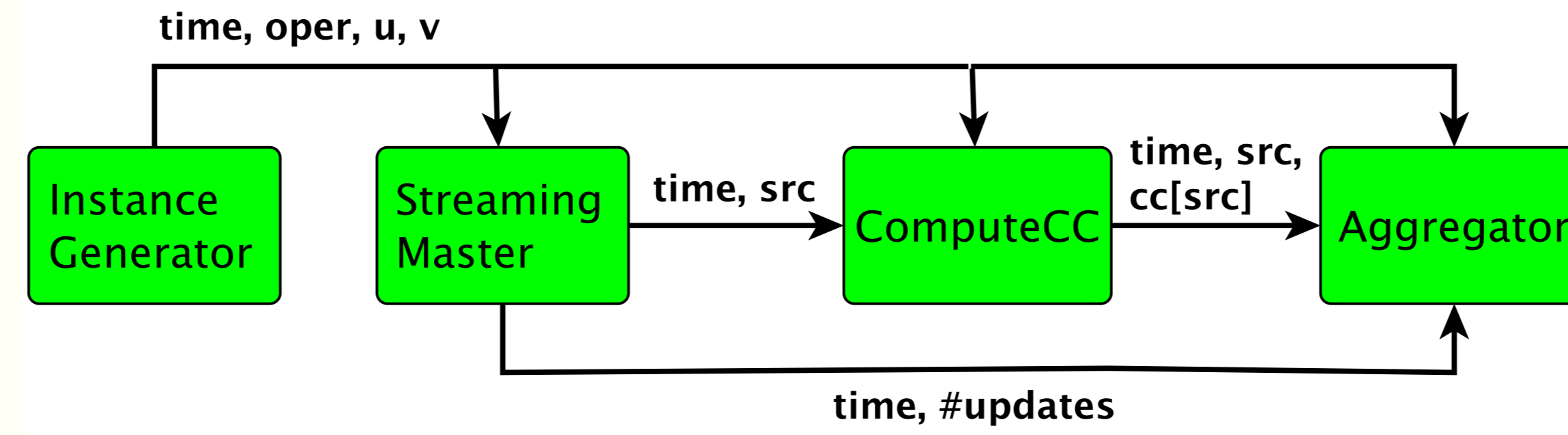
## Incremental Algorithm

- Upon insertion of  $(u, v)$
- Ignore temporarily identical vertices
- Consider  $(u, v)$ 's biconnected component
- Compute BFS from  $u$  and  $v$  (before edge insertion)
- Flag all  $s \neq u, v$ , s.t.  $|d(u, s) - d(v, s)| > 1$ ,
- Add  $(u, v)$  to the graph
- Compute  $cc[s]$  for all flagged  $s$
- Change  $cc[s]$  in other biconnected components
- Copy  $cc[s]$  for identical vertices

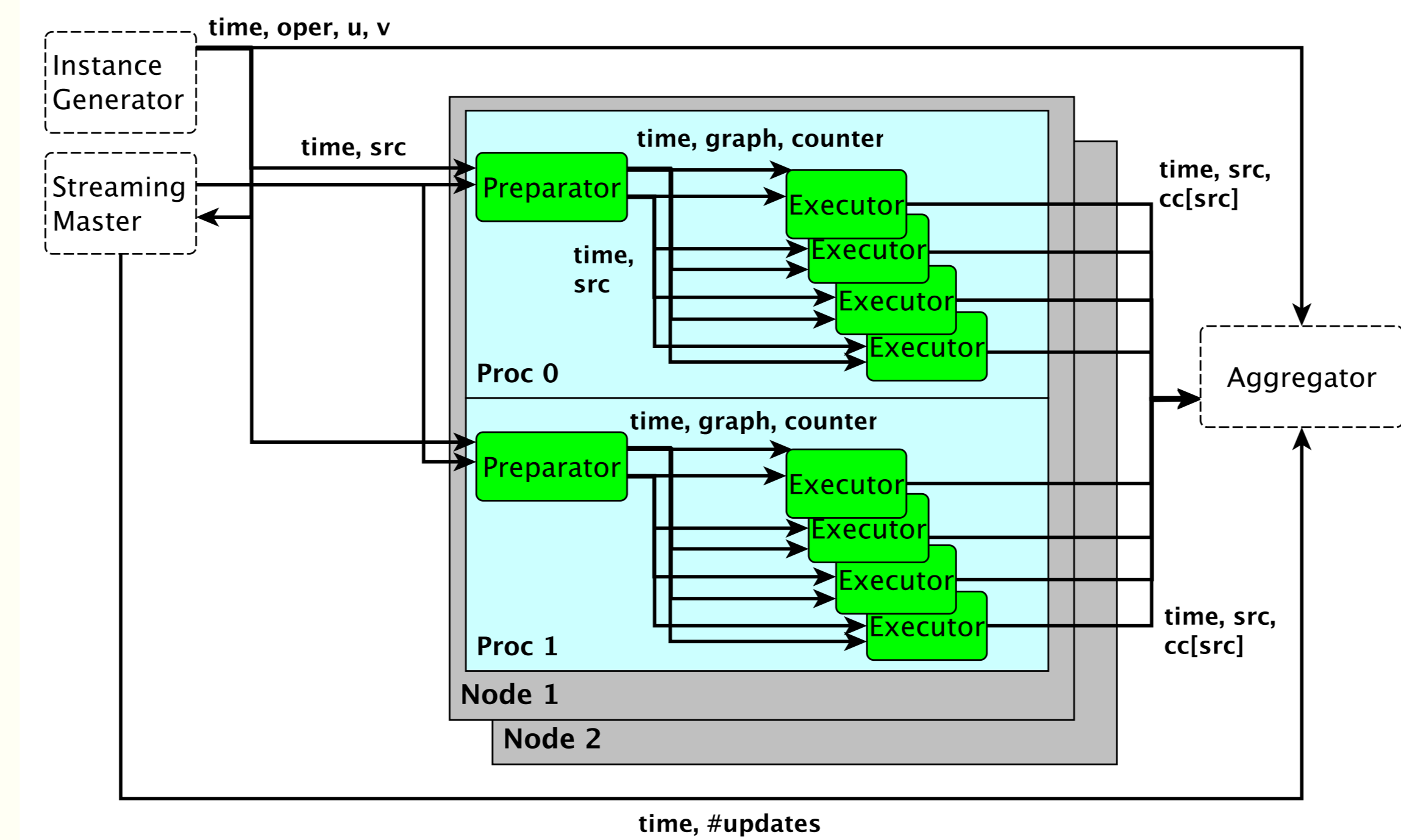
## Acknowledgment

This work was partially supported by the Defense Threat Reduction Agency grant HDTRA1-14-C-0007.

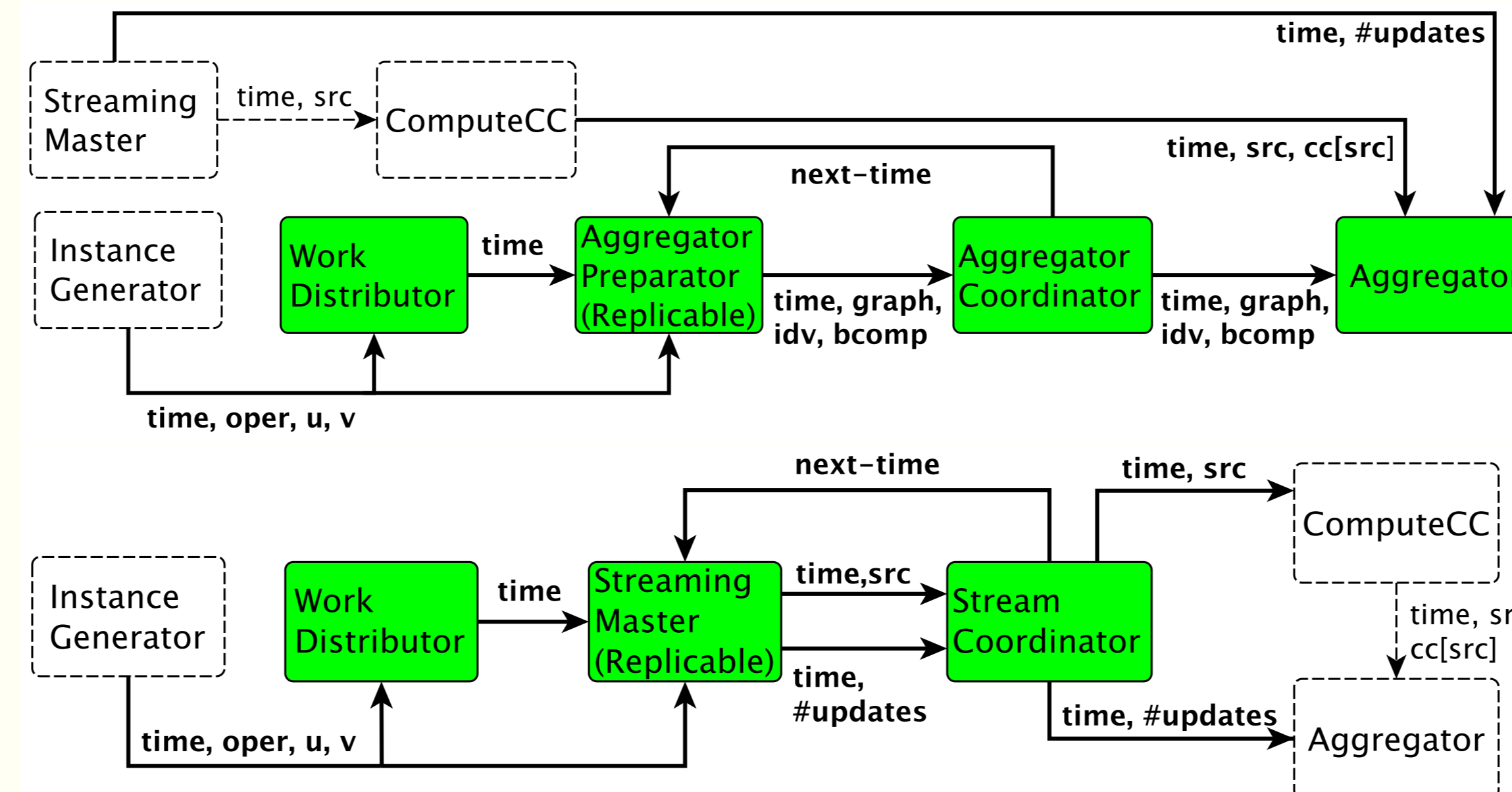
## Streamer Layout



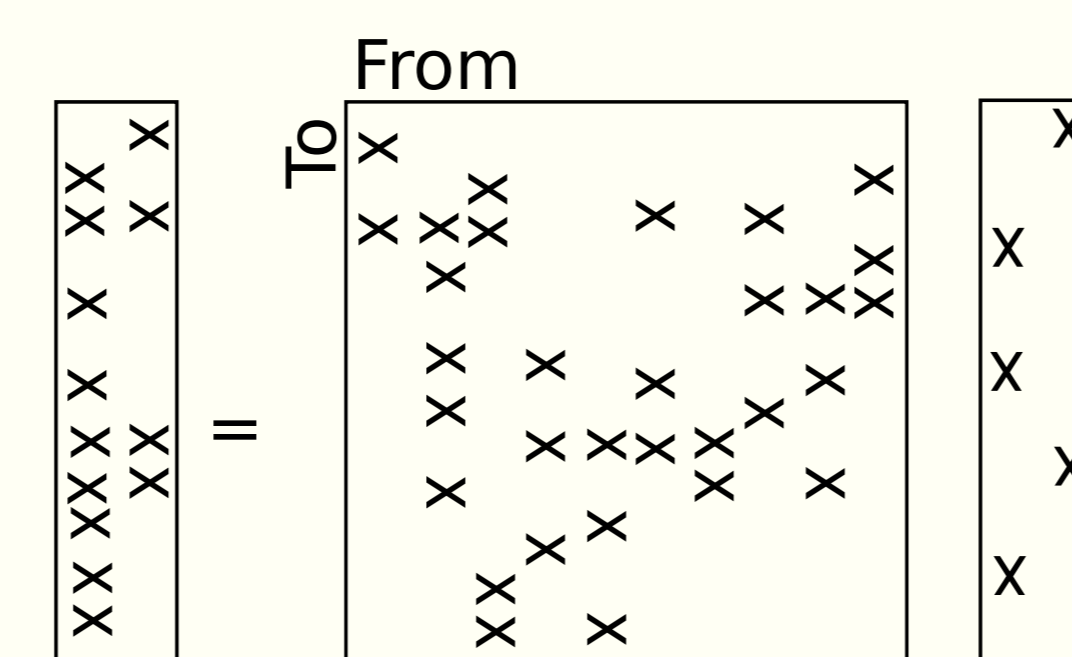
## Leveraging the NUMA Hierarchy



## Making Aggregator and StreamingMaster Replicable



## Vectorized Bottom-Up Traversal

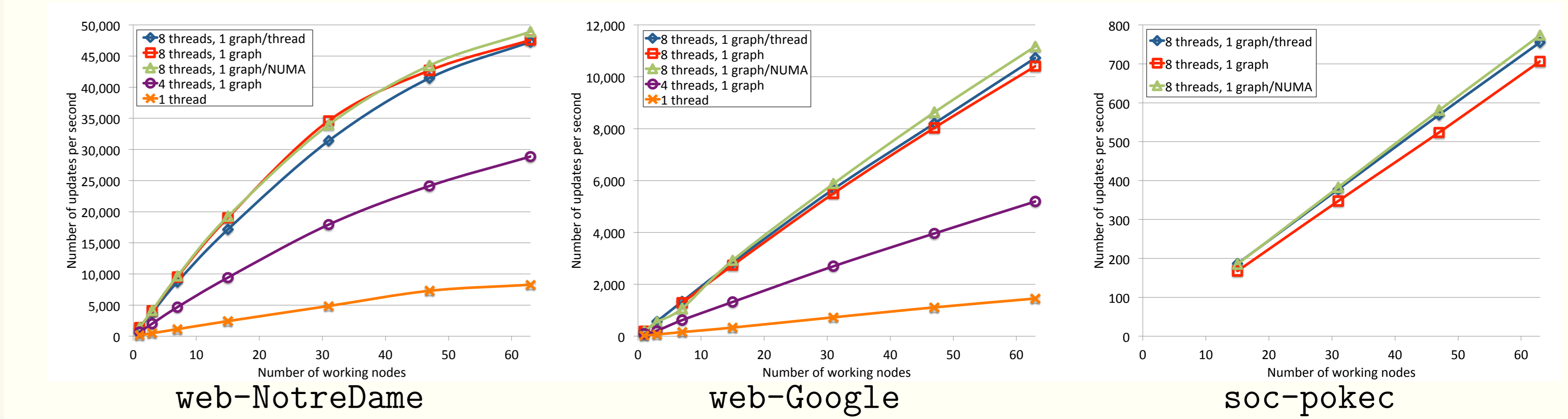


Multiple traversals can be executed simultaneously to leverage the large vector registers of modern processors.

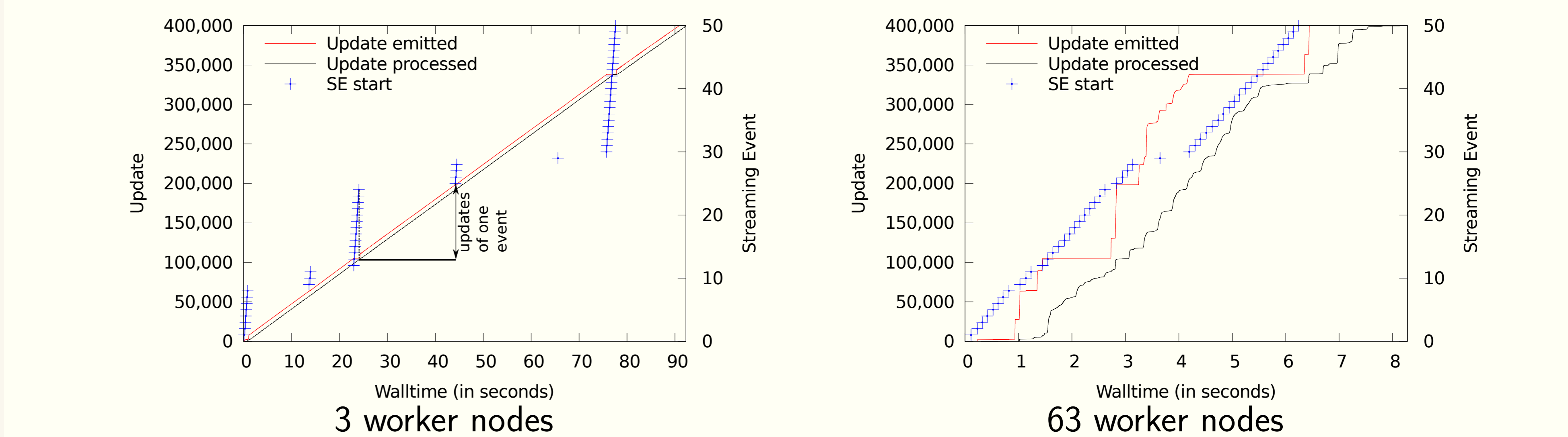
## Test Machine

- 64 nodes computational cluster:
  - 2 Intel Xeon E5520 @2.27Ghz
  - 4 cores (8 cores per node)
  - 2-way SMT (16 threads per node)
  - 8MB L3 cache (shared per proc)
  - 48 GB of main memory
  - 20 Gbps InfiniBand

## Experimental Results



## Log Analysis of web-NotreDame



## Using Replicable StreamingMaster and Aggregator (and Vectorization)

